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U.S. Application No. 09/940,985

Amendments to the Claims:

This listing of claims will replace all prior versions, and listings, of claims in the application:

Listing of Claims:

1. (Withdrawn) A tamper-resistant processing method comprising the steps of:
 - (1) storing a secret key index x corresponding to a public key of RSA (e , N ; modulus N being a product of 2 primes p and q) in a storage device;
 - (2) inputting a ciphertext Y through an input means;
 - (3) calculating yp , a remainder of y , based on a modulus of either P or its multiple and yq , a remainder of y , based on a modulus of either Q or its multiple; and
 - (4) when calculating Cp which is a remainder of yp^{xp} based on a modulus of either of p or its multiple and calculating Cq which is a remainder of yq^{xq} based on a modulus of either q or its multiple, where a remainder of x based on a modulus of either of $p-1$ or its multiple is put as xp , and a remainder of x based on a modulus of either $q-1$ or its multiple is put as xq ,
 - (4a) deciding which process (4b) or (4c) is to be executed for each processing of a bit block which is a bit string of at least 1 bit composing xp , xq ;
 - (4b) executing a predetermined modular exponentiation calculation on said bit block to be processed by xp and for storing the calculation result in the storage device;

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(4c) executing a predetermined modular exponentiation calculation on said bit block to be processed by xq and for storing the calculation result in the storage device;

(5) calculating RSA decryption calculation, $y^x \bmod N$ based on a difference between Cp and Cq , when the calculation of Cp about the whole of xp , and the calculation of Cq about the whole of xq are finished; and

(6) outputting the result of said RSA decryption calculation.

2. (Withdrawn) A tamper-resistant processing method of claim 1 wherein for said yp , yq , xp and xq , calculation be made as: $yp = y \bmod p$, $yq = y \bmod q$, $xp = x \bmod (p-1)$, $xq = x \bmod (q-1)$.

3. (Withdrawn) A tamper-resistant processing method of claim 1 wherein which one of said steps (4b) and (4c) is to be processed is determined with the use of a generated random number.

4. (Withdrawn) A tamper-resistant processing method of claim 1 wherein the process of said step (4a) is applied to a part of bit patterns of said xp or xq , and for a remaining part of the bit patterns, after said either one of step (4b) or (4c) is processed, another one is processed.

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5. (Currently amended) A tamper-resistant processing method comprising the steps of:

(a) deciding which step is to be selected out of the following steps (b) and (c) for each processing of one operation unit;

(b) after transferring one operation unit in the bit pattern of data A in a memory in order of ~~its bit-bit~~ sequence of said data A to a first register R1, transferring one operation unit in the bit pattern of data B in the memory in order of ~~its bit-bit~~ sequence of said data B to a second register R2;

(c) after transferring one operation unit in the bit pattern of said data B in order of ~~its bit-bit~~ sequence of said data B to said second register R2, transferring one operation unit in the bit pattern in said data A in order of ~~its bit-bit~~ sequence of said data A to said first register R1;

(d) executing a predetermined arithmetic operation on the contents of said first register R1 and the contents of said second register R2;

(e) storing the result of said arithmetic operation in the memory,

(f) repeating the steps from (a) through (e) until said arithmetic operation for said data A and said data B is finished.

6. (Currently amended) A tamper-resistant processing method comprising the steps of:

(a) deciding which step is to be selected out of the following steps (b) and (c) for each processing of one operation unit;

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(b) after transferring one operation unit of data A in a memory in order of ~~its bit~~ bit sequence of said data A to a first register R1, transferring one operation unit of data B in the memory in order of ~~its bit~~ bit sequence of said data B to a second register R2;

(c) after transferring said one operation unit of the data A in order of ~~its bit~~ bit sequence of said data A to said second register R2, transferring said one operation unit of the data B in order of ~~its bit~~ bit sequence of said data B to said first register R1;

(d) executing a predetermined arithmetic operation on the contents of said first register R1 and on the contents of said second register R2;

(e) storing the result of said arithmetic operation in the memory;

(f) repeating the steps from (a) through (e) until said arithmetic operation on said data A and said data B is finished.

7. (Previously presented) A tamper-resistant processing method of claim 6 wherein which one out of said steps (b) and (c) is to be processed is determined with the use of a generated random number.

8. (Original) A tamper-resistant processing method of claim 6 wherein said predetermined arithmetic operation is the operation for an arithmetic sum.

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9. (Original) A tamper-resistant processing method of claim 6 wherein said predetermined arithmetic operation is the operation for an arithmetic product.

10. (Original) A tamper-resistant processing method of claim 6 wherein said predetermined arithmetic operation is any one of the logical sum OR, logical product AND, and exclusive logical sum EXOR.

11. (Currently amended) A tamper-resistant processing method comprising the steps of:

- (a) randomly selecting any one of unprocessed one operation unit in a bit pattern of data A in a memory;
- (b) transferring said one operation unit of said data A selected to a first register R1;
- (c) transferring one operation unit in a bit pattern of data B in the memory corresponding to said one operation unit of said data A selected to a second register R2;
- (d) executing a predetermined arithmetic operation for the contents of said first register R1 and the contents of said second register R2;
- (e) storing a result of said arithmetic operation in the memory;
- (f) repeating the steps from ~~(4)~~(a) through ~~(5)~~(e) until said arithmetic operation is finished on said data A and said data B, wherein neither said data A nor said data B are randomly arranged bits.

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12. (Original) A tamper-resistant processing method of claim 11 wherein corresponding to a generated random number, said unprocessed one operation unit is selected.

13. (Original) A tamper-resistant processing method of claim 11 wherein said predetermined arithmetic operation is any one of logical sum OR, logical product AND, and exclusive logical sum EXOR.